Learning Objectives

At the end of the class you should be able to:

- explain how a generic searching algorithm works
- demonstrate how depth-first search will work on a graph
- demonstrate how breadth-first search will work on a graph
- predict the space and time requirements for depth-first and breadth-first searches



• Generic search algorithm: given a graph, start nodes, and goal nodes, incrementally explore paths from the start nodes.

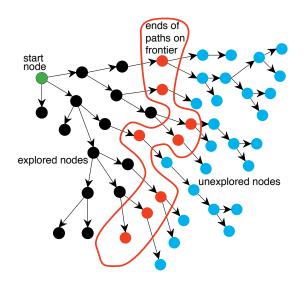


- Generic search algorithm: given a graph, start nodes, and goal nodes, incrementally explore paths from the start nodes.
- Maintain a frontier of paths from the start node that have been explored.

- Generic search algorithm: given a graph, start nodes, and goal nodes, incrementally explore paths from the start nodes.
- Maintain a frontier of paths from the start node that have been explored.
- As search proceeds, the frontier expands into the unexplored nodes until a goal node is encountered.

- Generic search algorithm: given a graph, start nodes, and goal nodes, incrementally explore paths from the start nodes.
- Maintain a frontier of paths from the start node that have been explored.
- As search proceeds, the frontier expands into the unexplored nodes until a goal node is encountered.
- The way in which the frontier is expanded defines the search strategy.

Problem Solving by Graph Searching





Graph Search Algorithm

```
Input: a graph,
         a set of start nodes.
          Boolean procedure goal(n) that tests if n is a goal node.
frontier := \{\langle s \rangle : s \text{ is a start node}\}
while frontier is not empty:
         select and remove path \langle n_0, \ldots, n_k \rangle from frontier
         if goal(n_k)
            return \langle n_0, \ldots, n_k \rangle
         for every neighbor n of n_k
            add \langle n_0, \ldots, n_k, n \rangle to frontier
end while
```

Graph Search Algorithm

- Which value is selected from the frontier at each stage defines the search strategy.
- The neighbors define the graph.
- goal defines what is a solution.
- If more than one answer is required, the search can continue from the return.

Optimality Criteria

- Often we don't want any solution, but the best solution or optimal solution.
- Costs on arcs give costs on paths. We want the least-cost path to a goal.



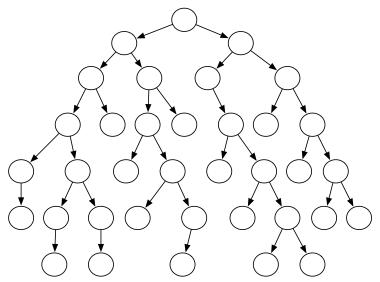
Depth-first Search

- Depth-first search treats the frontier as a stack
- It always selects one of the last elements added to the frontier.
- If the list of paths on the frontier is $[p_1, p_2, ...]$
 - p_1 is selected. Paths that extend p_1 are added to the front of the stack (in front of p_2).
 - \triangleright p_2 is only selected when all paths from p_1 have been explored.

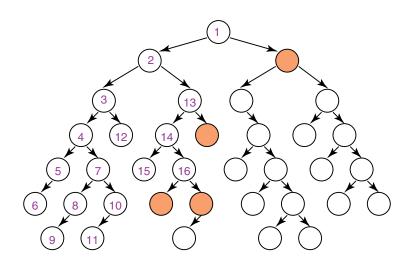


Illustrative Graph - Depth-first search

Start node is at the top.

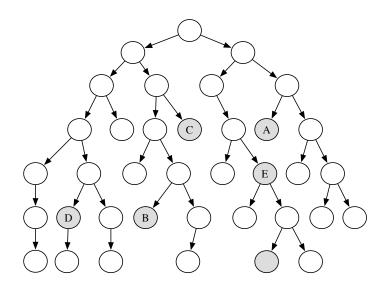


Illustrative Graph — Depth-first Search





Which shaded goal will depth-first search find first?





 Does depth-first search guarantee to find the path with fewest arcs?



- Does depth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?



- Does depth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?
- What is the time complexity as a function of length of the path selected?



- Does depth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?
- What is the time complexity as a function of length of the path selected?
- What is the space complexity as a function of length of the path selected?



- Does depth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?
- What is the time complexity as a function of length of the path selected?
- What is the space complexity as a function of length of the path selected?
- How does the goal affect the search?



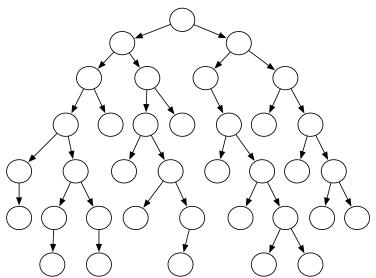
Breadth-first Search

- Breadth-first search treats the frontier as a queue.
- It always selects one of the earliest elements added to the frontier.
- If the list of paths on the frontier is $[p_1, p_2, ..., p_r]$:
 - p_1 is selected. Its neighbors are added to the end of the queue, after p_r .
 - \triangleright p_2 is selected next.

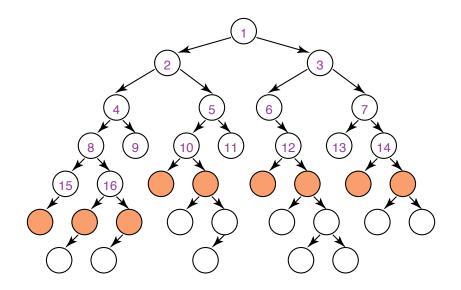


Illustrative Graph - Breadth-first search

Start node is at the top.

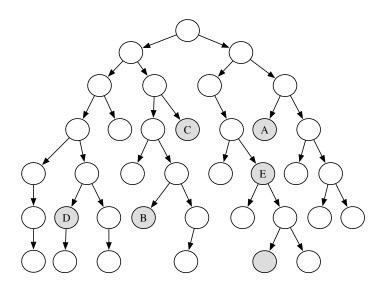


Illustrative Graph — Breadth-first Search





Which shaded goal will breadth-first search find first?



 Does breadth-first search guarantee to find the path with fewest arcs?



- Does breadth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?



- Does breadth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?
- What is the time complexity as a function of the length of the path selected?

- Does breadth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?
- What is the time complexity as a function of the length of the path selected?
- What is the space complexity as a function of the length of the path selected?

- Does breadth-first search guarantee to find the path with fewest arcs?
- What happens on infinite graphs or on graphs with cycles if there is a solution?
- What is the time complexity as a function of the length of the path selected?
- What is the space complexity as a function of the length of the path selected?
- How does the goal affect the search?



 Sometimes there are costs associated with arcs. The cost of a path is the sum of the costs of its arcs.

$$cost(\langle n_0,\ldots,n_k\rangle) = \sum_{i=1}^k cost(\langle n_{i-1},n_i\rangle)$$



 Sometimes there are costs associated with arcs. The cost of a path is the sum of the costs of its arcs.

$$cost(\langle n_0, \ldots, n_k \rangle) = \sum_{i=1}^k cost(\langle n_{i-1}, n_i \rangle)$$

- At each stage, lowest-cost-first search selects a path on the frontier with lowest cost.
- The frontier is a priority queue ordered by path cost.
- The first path to a goal is



 Sometimes there are costs associated with arcs. The cost of a path is the sum of the costs of its arcs.

$$cost(\langle n_0, \ldots, n_k \rangle) = \sum_{i=1}^k cost(\langle n_{i-1}, n_i \rangle)$$

- At each stage, lowest-cost-first search selects a path on the frontier with lowest cost.
- The frontier is a priority queue ordered by path cost.
- The first path to a goal is a least-cost path to a goal node.
- When arc costs are equal ⇒



 Sometimes there are costs associated with arcs. The cost of a path is the sum of the costs of its arcs.

$$cost(\langle n_0, \ldots, n_k \rangle) = \sum_{i=1}^k cost(\langle n_{i-1}, n_i \rangle)$$

- At each stage, lowest-cost-first search selects a path on the frontier with lowest cost.
- The frontier is a priority queue ordered by path cost.
- The first path to a goal is a least-cost path to a goal node.
- When arc costs are equal ⇒breadth-first search.



Summary of Search Strategies

| Strategy | Frontier Selection | Complete | Halts | Space |
|-------------------|--------------------|----------|-------|-------|
| Depth-first | Last node added | | | |
| Breadth-first | First node added | | | |
| Lowest-cost-first | Minimal $cost(p)$ | | | |

Complete — guaranteed to find a solution if there is one (for graphs with finite number of neighbours, even on infinite graphs)

Halts — on finite graph (perhaps with cycles).

Space — as a function of the length of current path



Summary of Search Strategies

| Strategy | Frontier Selection | Complete | Halts | Space |
|-------------------|--------------------|----------|-------|--------|
| Depth-first | Last node added | No | No | Linear |
| Breadth-first | First node added | Yes | No | Exp |
| Lowest-cost-first | Minimal $cost(p)$ | Yes | No | Exp |

Complete — guaranteed to find a solution if there is one (for graphs with finite number of neighbours, even on infinite graphs)

Halts — on finite graph (perhaps with cycles).

Space — as a function of the length of current path