

# Extensive Form Games and Backward Induction

ISCI 330 Lecture 13

February 27, 2007

# Lecture Overview

Recap

Subgame Perfection

Backward Induction

# Nash Equilibria

Given our new definition of pure strategy, we are able to reuse our old definitions of:

- ▶ mixed strategies
- ▶ best response
- ▶ Nash equilibrium

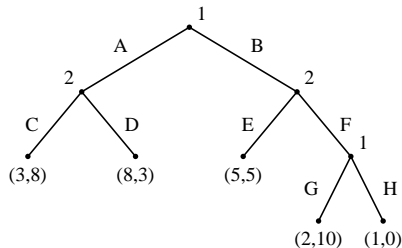
## Theorem

*Every perfect information game in extensive form has a PSNE*

This is easy to see, since the players move sequentially.

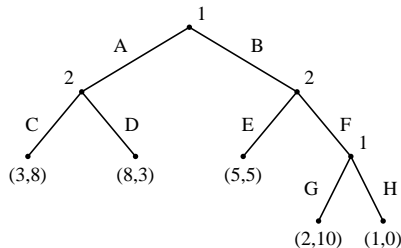
# Induced Normal Form

- ▶ In fact, the connection to the normal form is even tighter
  - ▶ we can “convert” an extensive-form game into normal form



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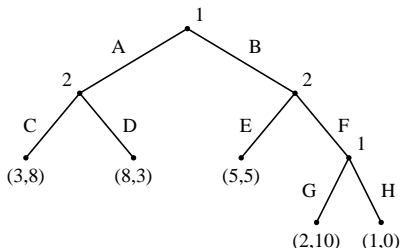
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<i>AG</i>	3, 8	3, 8	8, 3	8, 3
<i>AH</i>	3, 8	3, 8	8, 3	8, 3
<i>BG</i>	5, 5	2, 10	5, 5	2, 10
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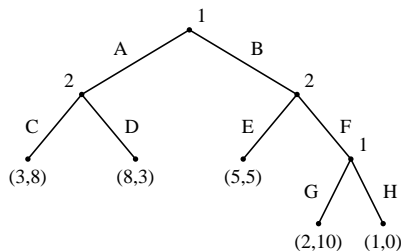


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- ▶ this illustrates the lack of compactness of the normal form
  - ▶ games aren't always this small
  - ▶ even here we write down 16 payoff pairs instead of 5

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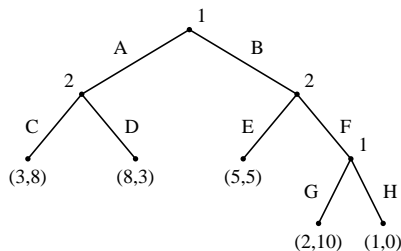


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- ▶ while we can write any extensive-form game as a NF, we can't do the reverse.
  - ▶ e.g., matching pennies cannot be written as a perfect-information extensive form game

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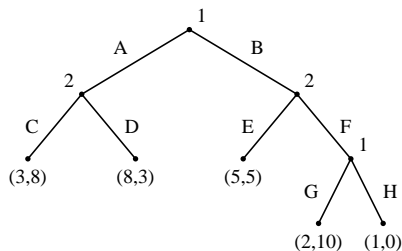
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- ▶ What are the (three) pure-strategy equilibria?



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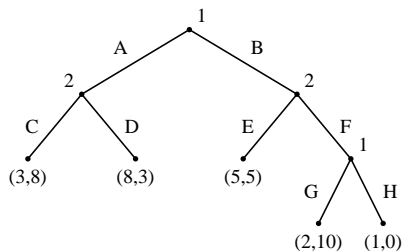


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- ▶ What are the (three) pure-strategy equilibria?
  - ▶  $(A, G), (C, F)$
  - ▶  $(A, H), (C, F)$
  - ▶  $(B, H), (C, E)$

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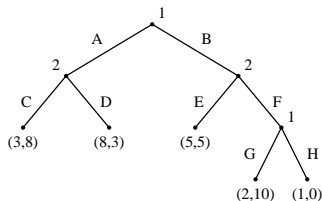
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Recap

Subgame Perfection

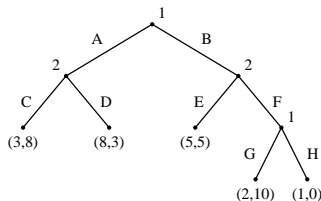
Backward Induction

# Subgame Perfection



- ▶ There's something intuitively wrong with the equilibrium  $(B, H), (C, E)$ 
  - ▶ Why would player 1 ever choose to play H if he got to the second choice node?
    - ▶ After all,  $G$  dominates  $H$  for him

# Subgame Perfection

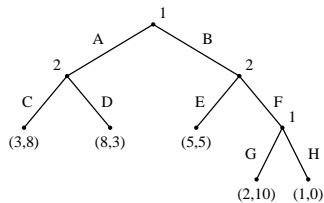


- ▶ There's something intuitively wrong with the equilibrium  $(B, H), (C, E)$ 
  - ▶ Why would player 1 ever choose to play H if he got to the second choice node?
    - ▶ After all,  $G$  dominates  $H$  for him
  - ▶ He does it to threaten player 2, to prevent him from choosing  $F$ , and so gets 5
    - ▶ However, this seems like a non-credible threat
    - ▶ If player 1 reached his second decision node, would he really follow through and play  $H$ ?

# Formal Definition

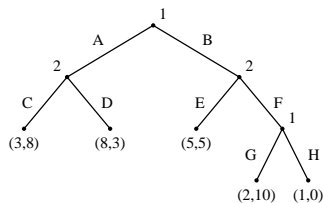
- ▶ Define **subgame of  $G$  rooted at  $h$** :
  - ▶ the restriction of  $G$  to the descendants of  $H$ .
- ▶ Define **set of subgames of  $G$** :
  - ▶ subgames of  $G$  rooted at nodes in  $G$
- ▶  $s$  is a **subgame perfect equilibrium** of  $G$  iff for any subgame  $G'$  of  $G$ , the restriction of  $s$  to  $G'$  is a Nash equilibrium of  $G'$
- ▶ Notes:
  - ▶ since  $G$  is its own subgame, every SPE is a NE.
  - ▶ this definition rules out “non-credible threats”

# Back to the Example



- ▶ Which equilibria from the example are subgame perfect?

# Back to the Example



- ▶ Which equilibria from the example are subgame perfect?
  - ▶  $(A, G), (C, F)$  is subgame perfect
  - ▶  $(B, H)$  is a non-credible threat, so  $(B, H), (C, E)$  is not subgame perfect
  - ▶  $(A, H)$  is also non-credible, even though  $H$  is “off-path”



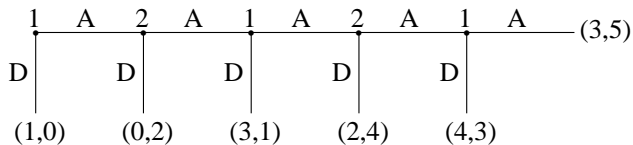
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# Centipede Game



- ▶ Play this as a fun game...