

Breadth-first Search; Search with Costs

CPSC 322 – Search 3

Textbook §3.5

Lecture Overview

1 Recap

2 Breadth-First Search

Graph Search Algorithm

Input: a graph,
a set of start nodes,
Boolean procedure $goal(n)$ that tests if n is a goal node.
 $frontier := \{\langle s \rangle : s \text{ is a start node}\};$
while $frontier$ is not empty:
 select and remove path $\langle n_0, \dots, n_k \rangle$ from $frontier$;
 if $goal(n_k)$
 return $\langle n_0, \dots, n_k \rangle$;
 for every neighbor n of n_k
 add $\langle n_0, \dots, n_k, n \rangle$ to $frontier$;
end while

- After the algorithm returns, it can be asked for more answers and the procedure continues.
- Which value is selected from the frontier defines the search strategy.
- The *neighbor* relationship defines the graph.
- The *goal* function defines what is a solution.

Depth-first Search

- **Depth-first search** treats the frontier as a stack
 - It always selects one of the last elements added to the frontier.

- **Complete** when the graph has no cycles and is finite
- **Time complexity** is $O(b^m)$
- **Space complexity** is $O(bm)$

DFS Example

- `http://aispace.org/search/`
- “simple tree graph”

Using Depth-First Search

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 - you know how to order nodes in the list of neighbours so that solutions will be found relatively quickly

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- When is DFS **inappropriate**?
 - some paths have infinite length
 - the graph contains cycles
 - some solutions are very deep, while others are very shallow

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Breadth-first Search

- Breadth-first search treats the frontier as a **queue**
 - it always selects one of the earliest elements added to the frontier.
- **Example:**
 - the frontier is $[p_1, p_2, \dots, p_r]$
 - neighbours of p_1 are $\{n_1, \dots, n_k\}$
- What happens?
 - p_1 is selected, and tested for being a goal.
 - Neighbours of p_1 follow p_r at the end of the frontier.
 - Thus, the frontier is now $[p_2, \dots, p_r, (p_1, n_1), \dots, (p_1, n_k)]$.
 - p_2 is selected next.

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 - The time complexity is $O(b^m)$: must examine every node in the tree.
 - The order in which we examine nodes (BFS or DFS) makes no difference to the worst case: search is unconstrained by the goal.

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- What is the **space complexity**?
 - Space complexity is $O(b^m)$: we must store the whole frontier in memory

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 - although all solutions may not be shallow, at least some are
 - there may be infinite paths

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- When is BFS **inappropriate**?
 - space is limited
 - all solutions tend to be located deep in the tree
 - the branching factor is very large