Reasoning Under Uncertainty: Variable elimination

Computer Science cpsc322, Lecture 30

(Textbook Chpt 6.4)

Nov, 20, 2012
Lecture Overview

• Recap Intro Variable Elimination
• Variable Elimination
  • Simplifications
  • Example
  • Independence
• Where are we?
Bnet Inference: General

• Suppose the variables of the belief network are $X_1, \ldots, X_n$.
• $Z$ is the query variable.
• $Y_1 = v_1, \ldots, Y_j = v_j$ are the observed variables (with their values).
• $Z_1, \ldots, Z_k$ are the remaining variables.

• What we want to compute: $P(Z \mid Y_1 = v_1, \ldots, Y_j = v_j)$.

• We can actually compute: $P(Z, Y_1 = v_1, \ldots, Y_j = v_j)$.

\[
P(Z \mid Y_1 = v_1, \ldots, Y_j = v_j) = \frac{P(Z, Y_1 = v_1, \ldots, Y_j = v_j)}{P(Y_1 = v_1, \ldots, Y_j = v_j)} = \frac{P(Z, Y_1 = v_1, \ldots, Y_j = v_j)}{\sum_Z P(Z, Y_1 = v_1, \ldots, Y_j = v_j)}
\]
Inference with Factors

We can compute $P(Z, Y_1=v_1, \ldots, Y_j=v_j)$ by

- expressing the joint as a factor,

$$f(Z, Y_1, \ldots, Y_j, Z_1, \ldots, Z_j)$$

- assigning $Y_1=v_1, \ldots, Y_j=v_j$

- and summing out the variables $Z_1, \ldots, Z_k$

$$P(Z, Y_1=v_1, \ldots, Y_j=v_j) = \sum_{Z_k} \cdots \sum_{Z_1} f(Z, Y_1, \ldots, Y_j, Z_1, \ldots, Z_k)_{Y_1=v_1, \ldots, Y_j=v_j}$$
Variable Elimination Intro (1)

\[ P(Z, Y_1 = v_1, \ldots, Y_j = v_j) = \sum_{Z_k} \cdots \sum_{Z_1} f(Z, Y_1, \ldots, Y_j, Z_1, \ldots, Z_k)_{Y_1 = v_1, \ldots, Y_j = v_j} \]

- Using the chain rule and the definition of a Bnet, we can write \( P(X_1, \ldots, X_n) \) as
  \[ \prod_{i=1}^n P(X_i \mid pX_i) \]
- We can express the joint factor as a product of factors
  \[ f(Z, Y_1, \ldots, Y_j, Z_1, \ldots, Z_j) \prod_{i=1}^n f(X_i, pX_i) \]
Variable Elimination Intro (2)

Inference in belief networks thus reduces to computing “the sums of products….”

\[ P(Z, Y_1 = v_1, \ldots, Y_j = v_j) = \sum_{Z_k} \cdots \sum_{Z_1} \prod_{i=1}^{n} f(X_i, pX_i) \]

1. Construct a factor for each conditional probability.
2. In each factor assign the observed variables to their observed values.
3. Multiply the factors
4. For each of the other variables \( Z_i \in \{Z_1, \ldots, Z_k\} \), sum out \( Z_i \)
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How to simplify the Computation?

- Assume we have turned the CPTs into factors and performed the assignments

\[
\sum_{Z_k} \ldots \sum_{Z_1} \prod_{i=1}^{n} f(X_i, pX_i)_{Y_1=v_1, \ldots, Y_j=v_j} \rightarrow \sum_{Z_k} \ldots \sum_{Z_1} \prod_{i=1}^{n} f(\text{vars}X_i)
\]

Let’s focus on the basic case, for instance...

\[
\sum_{A} f(C, D) \times f(A, B, D) \times f(E, A) \times f(D)
\]
How to simplify: basic case

Let’s focus on the basic case.  \[
\sum_{Z_1} \prod_{i=1}^{n} f(\text{vars}X_i) = \sum_{A} f(C, D) \times f(A, B, D) \times f(E, A) \times f(D)
\]

• How can we compute efficiently?

Factor out those terms that don't involve \(Z_1\)!

\[
\left( \prod_{i : Z_1 \notin \text{vars}X_i} f(\text{vars}X_i) \right) \times \left( \sum_{Z_1} \prod_{i : Z_1 \in \text{vars}X_i} f(\text{vars}X_i) \right)
\]

Do not contain \(Z_1\)  
Do contain

\[
+(C, D) \times +(D) \times \sum_{A} +(A, B, D) \times +(E, A)
\]
General case: Summing out variables efficiently

\[ \sum_{Z_k} \cdots \sum_{Z_1} f_1 \times \cdots \times f_h = \sum_{Z_k} \cdots \sum_{Z_2} (f_1 \times \cdots \times f_i) \left( \sum_{Z_1} f_{i+1} \times \cdots \times f_h \right) \]

Now to sum out a variable \( Z_2 \) from a product \( f_1 \times \cdots \times f_i \times f' \) of factors, again partition the factors into two sets:

- F: those that contain \( Z_2 \)
- F: those that do not

\[ \prod_{F} \times \sum_{Z_2} \prod_{F} \]
Analogy with “Computing sums of products”

This simplification is similar to what you can do in basic algebra with *multiplication* and *addition*

- It takes 14 multiplications or additions to evaluate the expression

\[ ab + ac + ad + aeh + afh + agh. \]

- This expression be evaluated more efficiently….

\[ a(b + c + d) + h(e + f + g) \]
Is there only one way to simplify?

\[ P(G, D=t) = \sum_{A,B,C} f(A, G) f(B, A) f(C, G) f(B, C) \]

\[ P(G, D=t) = \sum_A f(A, G) \sum_B f(B, A) \sum_C f(C, G) f(B, C) \]

\[ P(G, D=t) = \sum_A f(A, G) \sum_C f(C, G) \sum_B f(B, C) f(B, A) \]
Variable elimination algorithm: Summary

\[ P(Z \mid Y_1, \ldots, Y_j, Z_1, \ldots, Z_j) \]

To compute \( P(Z \mid Y_1=v_1, \ldots, Y_j=v_j) \):

1. Construct a factor for each conditional probability.
2. Set the observed variables to their observed values.
3. Given an elimination ordering, simplify/decompose sum of products
4. Perform products and sum out \( Z_i \)
5. Multiply the remaining factors (all in ?)
6. Normalize: divide the resulting factor \( f(Z) \) by \( \sum_Z f(Z) \).

A. \( Y_1=v_1 \)  B. \( Y_2 \)
C. \( Z_2 \)  D. \( Z \)
Variable elimination algorithm: Summary

\[ P(Z\mid Y_1, \ldots, Y_j, Z_1, \ldots, Z_j) \]

To compute \( P(Z\mid Y_1=v_1, \ldots, Y_j=v_j) \):

1. Construct a factor for each conditional probability.
2. Set the observed variables to their observed values.
3. Given an elimination ordering, simplify/decompose sum of products
4. Perform products and sum out \( Z_i \)
5. Multiply the remaining factors (all in \( f(Z) \))
6. Normalize: divide the resulting factor \( f(Z) \) by \( \sum_Z f(Z) \).
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• Where are we?
Variable elimination example

Compute \( P(G \mid H=h_1) \).

- \( P(G,H) = \sum_{A,B,C,D,E,F,I} P(A,B,C,D,E,F,G,H,I) \)
Variable elimination example

Compute $P(G \mid H=h_1)$.

- $P(G,H) = \sum_{A,B,C,D,E,F,I} P(A,B,C,D,E,F,G,H,I)$

Chain Rule + Conditional Independence:

$P(G,H) = \sum_{A,B,C,D,E,F,I} P(A)P(B\mid A)P(C)P(D\mid B,C)P(E\mid C)P(F\mid D)P(G\mid F,E)P(H\mid G)P(I\mid G)$
Variable elimination example (step1)

Compute $P(G \mid H=h_1)$.

- $P(G,H) = \sum_{A,B,C,D,E,F,I} P(A)P(B\mid A)P(C)P(D\mid B,C)P(E\mid C)P(F\mid D)P(G\mid F,E)P(H\mid G)P(I\mid G)$

Factorized Representation:

$P(G,H) = \sum_{A,B,C,D,E,F,I} f_0(A) f_1(B,A) f_2(C) f_3(D,B,C) f_4(E,C) f_5(F,D) f_6(G,F,E) f_7(H,G) f_8(I,G)$

- $f_0(A)$
- $f_1(B,A)$
- $f_2(C)$
- $f_3(D,B,C)$
- $f_4(E,C)$
- $f_5(F,D)$
- $f_6(G,F,E)$
- $f_7(H,G)$
- $f_8(I,G)$
Variable elimination example (step 2)

Compute $P(G \mid H=h_1)$.

Previous state:

$P(G,H) = \sum_{A,B,C,D,E,F,I} f_0(A) f_1(B,A) f_2(C) f_3(D,B,C) f_4(E,C) f_5(F, D) f_6(G,F,E) \ \underline{f_7(H,G)} \ f_8(I,G)$

Observe $H$:

$P(G,H=h_1) = \sum_{A,B,C,D,E,F,I} f_0(A) f_1(B,A) f_2(C) f_3(D,B,C) f_4(E,C) f_5(F, D) f_6(G,F,E) \ \underline{f_9(G)} \ f_8(I,G)$

- $f_0(A)$
- $f_1(B,A)$
- $f_2(C)$
- $f_3(D,B,C)$
- $f_4(E,C)$
- $f_5(F, D)$
- $f_6(G,F,E)$
- $\underline{f_7(H,G)}$
- $f_8(I,G)$

- $\underline{f_9(G)}$
Variable elimination example (steps 3-4)

Compute $P(G \mid H=h_1)$.

Previous state:

$$P(G, H) = \sum_{A,B,C,D,E,F,I} f_0(A) f_1(B,A) f_2(C) f_3(D,B,C) f_4(E,C) f_5(F, D) f_6(G,F,E) f_9(G) f_8(I,G)$$

Elimination ordering $A, C, E, I, B, D, F$:

$$P(G, H=h_1) = f_9(G) \sum_F \sum_D f_5(F, D) \sum_B \sum_I f_8(I,G) \sum_E f_6(G,F,E) \sum_C f_2(C) f_3(D,B,C) f_4(E,C) \sum_A f_0(A) f_1(B,A)$$

- $f_0(A)$
- $f_1(B,A)$
- $f_2(C)$
- $f_3(D,B,C)$
- $f_4(E,C)$
- $f_5(F, D)$
- $f_6(G,F,E)$
- $f_7(H,G)$
- $f_8(I,G)$
- $f_9(G)$
Variable elimination example (steps 3-4)

Compute \( P(G \mid H=h_1) \). Elimination ordering \( A, C, E, I, B, D, F \).

Previous state:
\[
P(G, H=h_1) = f_9(G) \sum_F \sum_D f_5(F, D) \sum_B \sum_I f_8(I, G) \sum_E f_6(G, F, E) \sum_C f_2(C) f_3(D, B, C) f_4(E, C) \sum_A f_0(A) f_1(B, A)
\]

Eliminate \( A \):
\[
P(G, H=h_1) = f_9(G) \sum_F \sum_D f_5(F, D) \sum_B f_{10}(B) \sum_I f_8(I, G) \sum_E f_6(G, F, E) \sum_C f_2(C) f_3(D, B, C) f_4(E, C)
\]
Variable elimination example (steps 3-4)

Compute $P(G \mid H=h_1)$. Elimination ordering $A, C, E, I, B, D, F$.

Previous state:

$$P(G,H=h_1) = \sum_f \sum_D f_5(F, D) \sum_B f_{10}(B) \sum_I f_8(I,G) \sum_E f_6(G,F,E) \sum_C f_2(C) f_3(D,B,C) f_4(E,C)$$

Eliminate $C$:

$$P(G,H=h_1) = \sum_f \sum_D f_5(F, D) \sum_B f_{10}(B) \sum_I f_8(I,G) \sum_E f_6(G,F,E) f_{12}(B,D,E)$$
Variable elimination example (steps 3-4)

Compute $P(G \mid H=h_1)$. Elimination ordering $A, C, E, I, B, D, F$.

Previous state:

\[
P(G, H=h_1) = f_9(G) \sum_F \sum_D f_5(F, D) \sum_B f_{10}(B) \sum_I f_8(I, G) \sum_E f_6(G, F, E) f_{12}(B, D, E)
\]

Eliminate E:

\[
P(G, H=h_1) = f_9(G) \sum_F \sum_D f_5(F, D) \sum_B f_{10}(B) f_{13}(B, D, F, G) \sum_I f_8(I, G)
\]

- $f_0(A)$
- $f_1(B, A)$
- $f_2(C)$
- $f_3(D, B, C)$
- $f_4(E, C)$
- $f_5(F, D)$
- $f_6(G, F, E)$
- $f_7(H, G)$
- $f_8(I, G)$
- $f_9(G)$
- $f_{10}(B)$
- $f_{12}(B, D, E)$
- $f_{13}(B, D, F, G)$
Variable elimination example (steps 3-4)

Compute $P(G \mid H=h_1)$. Elimination ordering $A$, $C$, $E$, $I$, $B$, $D$, $F$.

Previous state: $P(G, H=h_1) = f_9(G) \sum_F \sum_D f_5(F, D) \sum_B f_{10}(B) f_{13}(B, D, F, G) \sum_I f_8(I, G)$

Eliminate $I$:

$P(G, H=h_1) = f_9(G) f_{14}(G) \sum_F \sum_D f_5(F, D) \sum_B f_{10}(B) f_{13}(B, D, F, G)$

- $f_0(A)$
- $f_1(B, A)$
- $f_2(C)$
- $f_3(D, B, C)$
- $f_4(E, C)$
- $f_5(F, D)$
- $f_6(G, F, E)$
- $f_7(H, G)$
- $f_8(I, G)$
- $f_9(G)$
- $f_{10}(B)$
- $f_{12}(B, D, E)$
- $f_{13}(B, D, F, G)$
- $f_{14}(G)$
Variable elimination example (steps 3-4)

Compute $P(G | H=h_1)$. Elimination ordering $A, C, E, I, B, D, F$.

Previous state: $P(G,H=h_1) = f_9(G) f_{14}(G) \sum_F \sum_D f_5(F, D) \sum_B f_{10}(B) f_{13}(B,D,F,G)$

Eliminate B:

$P(G,H=h_1) = f_9(G) f_{14}(G) \sum_F \sum_D f_5(F, D) f_{15}(D,F,G)$
Variable elimination example (steps 3-4)

Compute \( P(G \mid H=h_1) \). Elimination ordering \( A, C, E, I, B, D, F \).

Previous state: \( P(G,H=h_1) = f_9(G) f_{14}(G) \sum_F \sum_D f_5(F, D) f_{15}(D,F,G) \)

Eliminate \( D \):
\[
P(G,H=h_1) = f_9(G) f_{14}(G) \sum_F f_{16}(F, G)
\]
Variable elimination example (steps 3-4)

Compute $P(G \mid H=h_1)$. Elimination ordering $A, C, E, I, B, D, F$.

Previous state: $P(G,H=h_1) = f_9(G) f_{14}(G) \sum_F f_{16}(F, G)$

Eliminate $F$:

$P(G,H=h_1) = f_9(G) f_{14}(G) f_{17}(G) \cdot f_9(G) \cdot f_{10}(B) \cdot f_{12}(B,D,E) \cdot f_{13}(B,D,F,G) \cdot f_{14}(G) \cdot f_{15}(D,F,G) \cdot f_{16}(F,G) \cdot f_{17}(G)$
Variable elimination example (step 5)

Compute $P(G \mid H=h_1)$. Elimination ordering $A, C, E, I, B, D, F$.

Previous state: $P(G,H=h_1) = f_9(G) f_{14}(G) f_{17}(G)$

Multiply remaining factors: 

$P(G,H=h_1) = f_{18}(G)$
Variable elimination example (step 6)

Compute $P(G \mid H=h_1)$. Elimination ordering $A, C, E, I, B, D, F$.

Previous state:

$P(G,H=h_1) = f_{18}(G)$

Normalize:

$$P(G \mid H=h_1) = \frac{f_{18}(G)}{\sum_{g \in \text{dom}(G)} f_{18}(G)}$$
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Complexity (not required)

- The complexity of the algorithm depends on a measure of complexity of the network.
- The size of a tabular representation of a factor is exponential in the number of variables in the factor.
- The **treewidth** of a network, given an elimination ordering, is the maximum number of variables in a factor created by summing out a variable, given the elimination ordering.
- The **treewidth** of a belief network is the minimum treewidth over all elimination orderings. The treewidth depends only on the graph structure and is a measure of the sparseness of the graph.
- The complexity of VE is exponential in the treewidth and linear in the number of variables.
- Finding the elimination ordering with minimum treewidth is NP-hard, but there is some good elimination ordering heuristics.
Variable elimination and conditional independence

• Variable Elimination looks incredibly painful for large graphs?
• We used conditional independence.....

\[ P(X_1, \ldots, X_n) = \prod P(X_i | \text{parents}(X_i)) \]

• Can we use it to make variable elimination simpler?

Yes, all the variables from which the query is conditional independent given the observations can be pruned from the Bnet
VE and conditional independence: Example

- All the variables from which the query is conditional independent given the observations can be pruned from the Bnet.

\[ P(G \mid H=v_1, F=v_2, C=v_3). \]

A. \( B, D, E \)  
B. \( E, D \)  
C. \( D, I \)  
D. \( B, D, A \)
VE and conditional independence: Example

- All the variables from which the query is conditional independent given the observations can be pruned from the Bnet.

\[ P(G \mid H=v_1, F=v_2, C=v_3). \]

G is conditionally independent from \( ABD \) given the observed vars \( H, F, C \).
VE and conditional independence: Example

- All the variables from which the query is conditional independent given the observations can be pruned from the Bnet.

\[ P(G \mid H = v_1, F = v_2, C = v_3) \]

- Both paths from G to D are blocked.

- G is conditionally independent from ABD given the observed vars H, F, I.
VE and conditional independence: Example

- All the variables from which the query is conditional independent given the observations can be pruned from the Bnet.

\[ P(G \mid H=v_1, F=v_2, C=v_3). \]

G is independent of DBA given C F.
Learning Goals for today’s class

You can:

• Carry out variable elimination by using factor representation and using the factor operations.

• Use techniques to simplify variable elimination.
Big Picture: R&R systems

Environment

Deterministic

Stochastic

Problem

Constraint Satisfaction

Query

Static

Sequential

Planning

Representation

Reasoning Technique

Deterministic

Arc Consistency

Search

SLS

Stochastic

Belief Nets

Var. Elimination

Decision Nets

Markov Processes

Value Iteration

Vars + Constraints

Search

Logics

STRIPS

Search

CSP

CSP

for CSP

for Inference

CSP

for complex planning

CSP

CSP

CPSC 322, Lecture 2
Answering Query under Uncertainty

- Probability Theory
- Static Belief Network & Variable Elimination
- Dynamic Bayesian Network
- Hidden Markov Models
- Monitoring (e.g., credit cards)
- BioInformatics
- Natural Language Processing
- Email spam filters
- Student Tracing in tutoring Systems

Some Applications:
- You will know a little
- Some Applications
Next Class

Probability and Time (TextBook 6.5)

Course Elements

• Work on Practice Exercise 6.C on variable elimination.

• Assignment 4 will be available tomorrow and due on Nov the 29th (last class).

• Fill out teaching evaluations. You should have received an email about this.